# Two combinatorial optimization problems coming from transportation

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CERMICS, Optimisation et Systmes

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First problem

## Joint aircraft-crew planning and stochastic shortest paths

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### Aircraft routing and crew pairings

PhD thesis AXEL PARMENTIER 2013-2016, AIR FRANCE



Objective: compute aircraft routes and crew routes minimizing costs

Graph D = (V, A): V =flights A = (u, v) in A if v can follow u



min 
$$\sum_{\substack{p \in \mathcal{P} \\ p \ni v}} c_p y_p$$
s.t. 
$$\sum_{\substack{p \ni v \\ y_p \in \{0, 1\}}} v \in V$$

$$\sum_{\substack{r \ni v \\ x_r \in \{0,1\}}} x_r = 1 \qquad v \in V$$

 $\sum x_r \leq \sum y_p \quad s \in S$ 

#### Algorithmic strategy

NP-hard problem, huge integer program : use of a mathheuristic.

Solve to optimality the crew pairing problem via column generation

$$\begin{array}{ll} \min & \sum_{p \in \mathcal{P}} c_p y_p \\ \text{s.t.} & \sum_{p \ni v} y_p = 1 \quad v \in V \\ & y_p \in \{0,1\} \quad p \in \mathcal{P} \end{array}$$

A posteriori check existence of a compatible aircraft routing solution via integer program

If not, solve again crew pairing via additional constraint

$$\sum_{\substack{r \ni v \\ r \ni s}} x_r = 1 \qquad v \in V$$
$$\sum_{\substack{r \ni s \\ x_r \in \{0, 1\}}} y_p \qquad s \in S, \ p \ni s$$

$$z_s \ge y_p$$
  $s \in S, \ p \ni s$   
 $\sum_{s \in \mathcal{F}} z_s \le |\mathcal{F}| - 1$ 

#### Delay

Control delay: working with  $p \in \mathcal{P}$  such that

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\mathbb{P}(\text{delay along } p > \tau) < \varepsilon
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"In 1 –  $\varepsilon$  of the cases, the delay is smaller than  $\tau$  minutes."

STOCHASTIC RESOURCE CONSTRAINT SHORTEST PATH as a subproblem:

Input. Graph D = (V, A), two vertices *s* and *t*, independent random travel times  $(X_a)_{a \in A}$ , costs  $(c_a)_{a \in A}$ .

Output. *s*-*t* path *P* with  $\mathbb{P}(\sum_{a \in P} X_a > \tau) < \varepsilon$  and with minimum  $\sum_{a \in P} c_a$ .

### Algorithm for Stochastic Resource Constraint Shortest Path

NP-hard problem. Idea: exhaustive and implicit enumeration of all paths.

Use stochastic lower bounds to discard uninteresting partial path.

Stochastic lower bound computed via a fixed-point equation:

$$\begin{cases} Z_t = 0 \\ Z_v = \bigwedge_{u \in N^+(v)} (X_{(v,u)} + Z_u) \quad v \in V \end{cases}$$

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Overall method currently implemented at AIR FRANCE.

Algorithm for Stochastic Resource Constraint Shortest Path able

- to solve instances with 1600 vertices and 6500 arcs in less than 15 seconds.
- to replace  $\mathbb{P}(\cdot > \tau)$  by other risk measures (such as CVaR).

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Second problem

### Shuttle scheduling

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#### Increase the capacity of the Chunnel

PhD thesis LAURENT DAUDET 2014-2017, Chaire EUROTUNNEL



Trains in the tunnel: Eurostars, freight trains, passagers shuttles (PAX), freight shuttles (HGV)

First objective: Increase capacity in HGV's Second objective: Increase capacity in PAX's

Constraints: Safety, "equity" schedule of Eurostars, fixed number of Eurostars and freight trains, fixed number of departures

#### Improvement?

EUROTUNNEL current rules: cyclic schedule, (one hour period), discretized (1 minute), 1 Eurostar every 30 minutes.

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Possible way to improve...

- "Relax" these rules to improve capacity.
- Decrease waiting time of passengers.

#### There is room for optimization

Currently: 10 shuttles per hour, in each direction.

Elementary experiments with integer programming (CPLEX) show

- with arbitrarily small discretization: 13 shuttles per hour
- with 1 Eurostar every 25-35 minutes: 11 shuttles per hour

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• with both + 2 hour cycle: 13.5 shuttles per hour.

### Decrease waiting time: simplified version

★ One type of mobile, one directions, known cumulated demand D(t) for all  $t \in [0, T]$ .

- \* Objective: minimize maximum waiting time of passengers.
- ★ Constraints:
  - Loading: constant rate ρ, starts at the arrival last passenger of shuttle.
  - Maximal load (*C*) of shuttles, fixed number of departures (*N*).

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• Safety constraints, speed limit.

Theorem Problem polynomially solvable.

#### A more realistic problem

★ One type of mobile, two directions, known cumulated demand D(t) for all  $t \in [0, T]$ .

\* Objective: minimize maximum waiting time of passengers.

- ★ Contraintes :
  - Loading: constant rate ρ, starts at the arrival last passenger of shuttle.
  - Maximal load (*C*) of shuttles, fixed number of departures (*N*).

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• Safety constraints, speed limit.

#### Proposed approach

#### Lagrangean heuristic.

- 1. remove the constraints with D(t) and add them to the objective function with a penalization
- 2. solve this problem (polynomial ~ simplified version)
- repeat with new penalization → convergence to a good lower bound
- 4. heuristically build a feasible solution from the lower bound

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Experiments currently carried out.

### Merci.

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